## **IxD Theory 2: Telecomunicazioni**

IUAV University of Venice

Visual and Multimedia Design graduate programme

**Programming the computer 2** 

© Gillian Crampton Smith + Philip Tabor 2009

## What does RobotProg teach us?

The basic elements of his program are:

do something commands conditional program flow loops

## What does RobotProg teach us?

Programming needs not only knowledge of the language, but also problem-solving strategy

We must make an abstraction of the real-world problem in order to manipulate it with the computer

A programming language has a limited number of commands that the computer can understand: we have to solve the problem using only these. Gillian Crampton Smith + Philip Tabor Telecomunicazioni

## Levels of programming languages

Visual programming languages – like Visual Basic or MaxSP

Meta-languages – that write code for you, like Dreamweaver, the web design program

High level languages – like C or Java or Processing

**Assembly language** – using mnemonics which are then translated into machine language

Machine language (using bit patterns) 0s and 1s

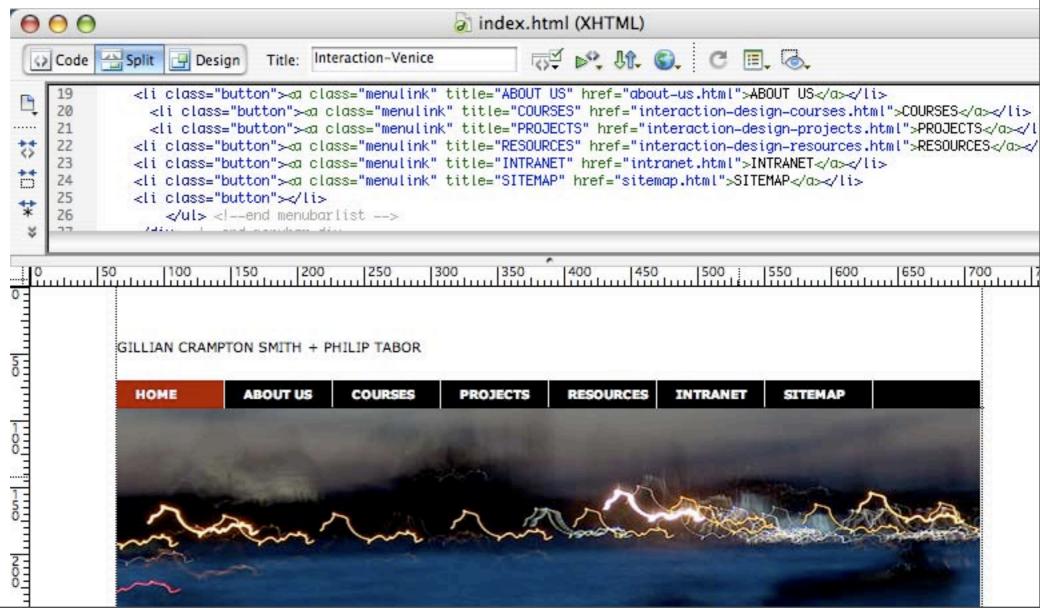
## Meta languages

## Meta languages

For example: Dreamweaver

- translated first into high-level language (e.g. Html) then into machine code by the computer

- even easier to understand BUT are less flexible.



## **High-level languages**

## High-level languages

For example: C++, Java, Processing, HTML and CSS

- are translated into machine code by the computer
- are easier to understand and avoid errors

We still need to know the right commands of our language, but we don't need to worry about which memory locations we are using

```
void setup()
{
size(200, 200);
background(0);
stroke(153);
loop();
}
int x = 0;
int y = 100;
int incx = 1;
int incy = 10;
void draw()
line(x, y, x+incx,y+incy);
x = x + incx;
y = y + incy;
if (x + y = 199) noLoop();
}
```

## MACHINE LANGUAGE AND ASSEMBLER

## How the computer manipulates 0s and 1s

Microprocessor:

**ALU:** (arithmetic logic unit) which does mathematical operations

**Registers:** temporary storage locations for operations in progress

Memory (RAM)

Longer-term storage locations

**Video memory** 

Screen buffers special memory mapped to the screen

## Machine language

Modern computers use 32-bit numbers to store machine language

## $000 \ 0001 \ 0010 \ 1011 \ 1000 \ 0000 \ 0010 \ 0000$

# Means: Add register 1 to register 2, and put the result in register 0.

AAARGH!!!!

Don't panic!

## Machine language

At the machine level, the program is written in patterns of 0s and 1s, each of which is a code for things like:

an instruction (e.g. ADD) a piece of data (e.g. 39) an address in memory to get some data an address in memory to put some data

Every microprocessor has a different instruction set

The first computers were programmed in machine language.



It was very slow and easy to make errors.

The Altair, an early desktop computer, had two rows of toggle switches: one for the data, and one for the address to put the data.

So assembly languages were invented which used 'mnemonics' (easy-to-remember codes). The computer then translated these codes into the 1s and Os of the machine

A line of assembly language has 4 elements:

Labels 'Opcodes' Operand(s) CommentsSUM:ADD\$t0, \$t1, \$t2;add 2 numbers;put the result in;register 2

It's difficult, but very fast.

Programs take very little memory so it's good for embedded microprocessors —in white goods, for instance: washing machines, cookers, vacuum cleaners, etc.

http://www.8052.com

- ACALL: Absolute Call
- ADD, ADDC: Add Accumulator (With Carry)
- <u>AJMP</u>: Absolute Jump
- ANL: Bitwise AND
- CJNE: Compare and Jump if Not Equal
- CLR: Clear Register
- CPL: Complement Register
- DA: Decimal Adjust
- DEC: Decrement Register
- DIV: Divide Accumulator by B
- DJNZ: Decrement Register and Jump if Not Zero
- INC: Increment Register
- JB: Jump if Bit Set
- JBC: Jump if Bit Set and Clear Bit
- JC: Jump if Carry Set
- JMP: Jump to Address
- JNB: Jump if Bit Not Set
- JNC: Jump if Carry Not Set
- JNZ: Jump if Accumulator Not Zero
- JZ: Jump if Accumulator Zero
- LCALL: Long Call
- LJMP: Long Jump
- MOV: Move Memory
- MOVC: Move Code Memory
- MOVX: Move Extended Memory
- MUL: Multiply Accumulator by B
- <u>NOP</u>: No Operation
- ORL: Bitwise OR
- POP: Pop Value From Stack
- <u>PUSH</u>: Push Value Onto Stack
- <u>RET</u>: Return From Subroutine
- <u>RETI</u>: Return From Interrupt
- <u>RL</u>: Rotate Accumulator Left
- <u>RLC</u>: Rotate Accumulator Left Through Carry
- <u>RR</u>: Rotate Accumulator Right
- <u>RRC</u>: Rotate Accumulator Right Through Carry
- SETB: Set Bit
- <u>SJMP</u>: Short Jump
- <u>SUBB</u>: Subtract From Accumulator With Borrow
- <u>SWAP</u>: Swap Accumulator Nibbles
- XCH: Exchange Bytes
- XCHD: Exchange Digits
- XRL: Bitwise Exclusive OR
- Undefined: Undefined Instruction

## Main Memory

#### **Microprocessor**

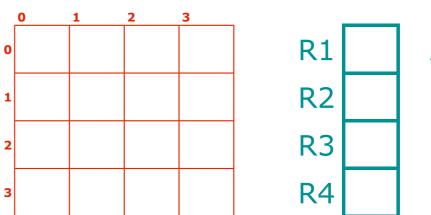
# micro-computer

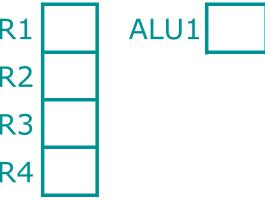
**The Terese** 

Our imaginary low-level computer has only 16 bytes of main memory and 100 bytes of video memory.

Its microprocessor has 4 registers and one arithmetic logic unit.

How would it work?





Remember that Asssembly language commands are like this:

SUM: ADD \$t0, \$t1, \$t2

So let's look at the command set for our imaginary computer. . . .

## **Terese micro-computer assember commands**

OPCODE	DE OPERATION		ON		
LR1	Load Register R1		operand 1 address or #number	operand 2 -	operand 3 -
LR2	Load Register R2		address or #number	-	-
C&M	Copy and move value in 2nd	А	address 1	address 2	
	address to 1st address	В	address 1	#number	
ADD	Add 2 values		1st register	2nd register	address to store the result
ЈМР	jump to a 'label'		label	-	-
JEQ	Checks if two values are equal AND jumps to a label		address 1	address 2	address or 'label' to jump to
INC	Increases a number in address 1 by the number in address 2	A	address 1	address 2	
END	Stops the program	В	address 1	#number	

## A Terese program

#### **Terese mini-computer: example of assembler program**

Label		opcode	operand 1	operand 2	operand 3
1	X:	C&M	0 0	#0	
2	Y:	C&M	01	#100	
3	PLUS X:	C&M	02	#1	
4	PLUS Y:	C&M	03	#10	
5	DRAW:	LR1	0 0		
6		LR2	01		
7		ADD	R1	<b>R2</b>	<b>R3</b>
8		C&M	<b>R3</b>	#1	
9		INC	X	02	
10		INC	У	03	
11		JEQ	#R3	199	END
12		JMP	DRAW		
13	END:	END			

#### Terese mini-computer: example of assembler program

	Label	орсо	de	operand 1	operand 2	operand 3	
1		X:	C&M	0 0	#0		; in memory location 0,0 put 0

Main Memory

2

3

ο X 1

0

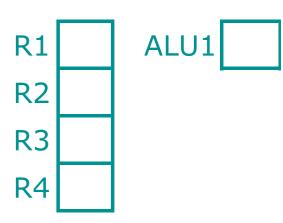
0

1

2

3

#### Microprocessor



Terese mini-computer: example of assembler program

	_abel	opcode	operand 1	operand 2	operand 3		Vide	eo M	1em	ory						
1	X:	C&M	0 0	#0	1	100	0 100	1 101	2 102	3 103	4 104	5 105	6 <b>106</b>	7 <b>107</b>	8 <b>108</b>	9 <b>109</b>
	n memor	y location	n 0,0 put t	the numb	er 0	; in r ; io r ;	memo ad R1 nory ad R2 nory 1 and t 1 (t ress i reme nory creme nory eck if video ed);	ory lo with 00 with 01 put r o ligh n R3 ent x 02 ent y 03 R3= men	catior catior conte conte conte esult t a pi with o 199 ( nory I	n 0,3 ents of ents of in R3 ixel) conte conte R3 is ocatio	put 1 of 144 in the nts of nts of	0 155 e		177	188	199

Main Memory

2

3

• X • Y

0

0

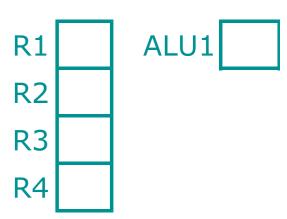
1

2

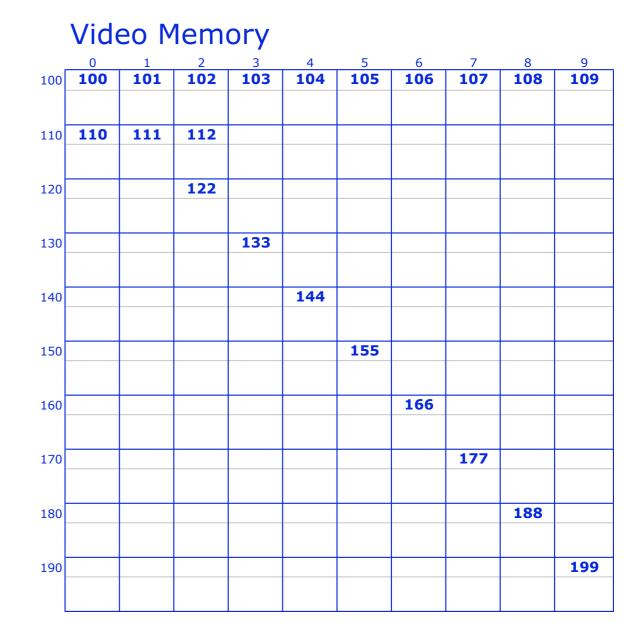
3

100

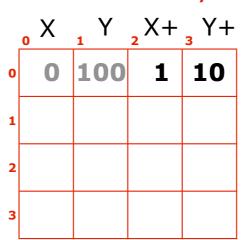
#### Microprocessor



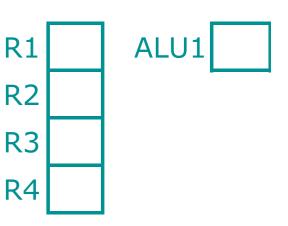
;in memory location 0,0 put 0 ;in memory location 0,1 put 100



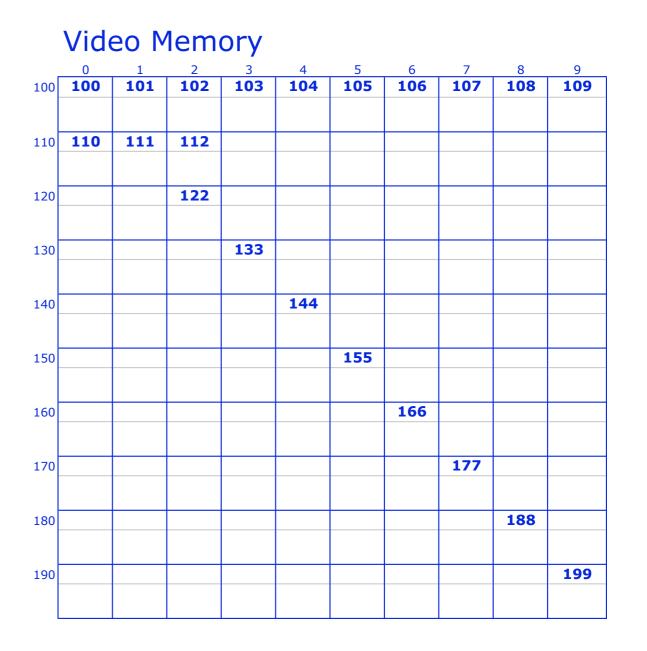
Main Memory



#### Microprocessor

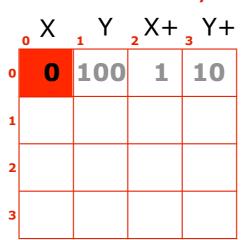


;in memory location 0,0 put 0
;in memory location 0,1 put 100
;in memory location 0,2 put 1
;in memory location 0,3 put 10

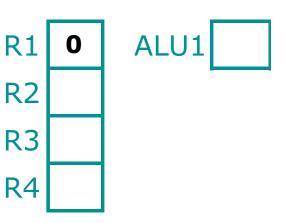


## **The Terese** micro-computer <u>X Y X+ Y+</u>

#### Main Memory

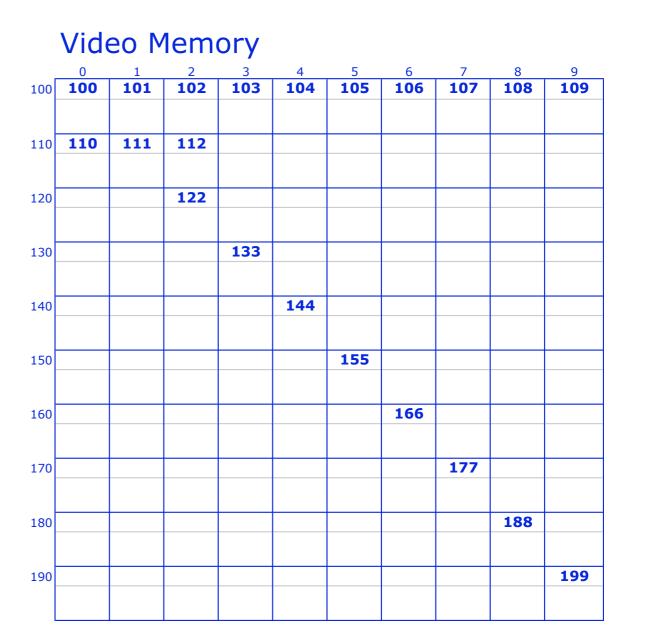


#### Microprocessor



; in memory location 0,0 put 0 ; in memory location 0,1 put 100 ; in memory location 0,2 put 1 ; in memory location 0,3 put 10

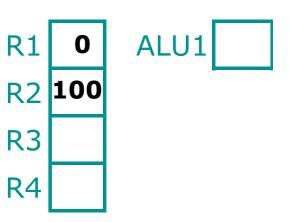
; load R1 with contents of memory 00



# The TereseMain Memorymicro-computerXY.XYY+

#### Main Memory <sup>0</sup> X <sup>1</sup> Y <sup>2</sup> X+ <sup>3</sup> Y+ **0 100 1 10** <sup>1</sup> 10 <sup>1</sup> 10 <sup>2</sup> 100 100 <sup>3</sup> 100 100 <sup>1</sup> 100 <sup>1</sup>

#### Microprocessor

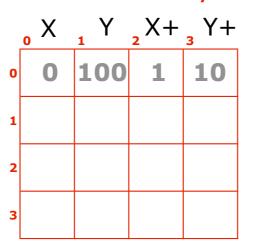


; in memory location 0,0 put 0 ; in memory location 0,1 put 100 ; in memory location 0,2 put 1 ; in memory location 0,3 put 10

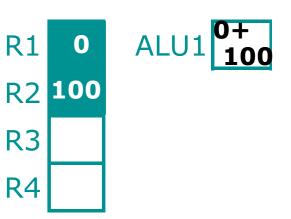
; load R1 with contents of memory 00

; load R2 with contents of memory 01

#### Main Memory



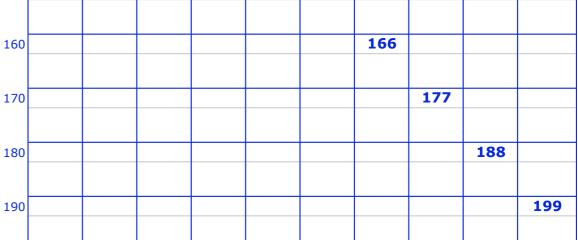
#### **Microprocessor**



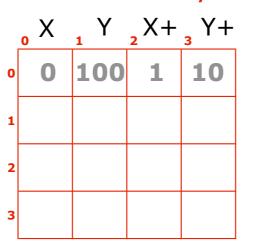
; in memory location 0,0 put 0 ; in memory location 0,1 put 100 ; in memory location 0,2 put 1 ; in memory location 0,3 put 10

# ; load R1 with contents of memory 00

; load R2 with contents of memory 01 ;add



#### Main Memory

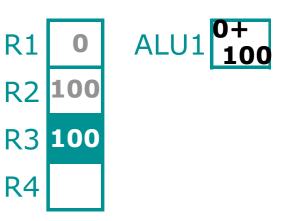


100

110

120

#### **Microprocessor**



8

108

9

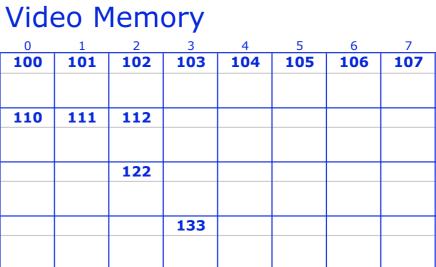
109

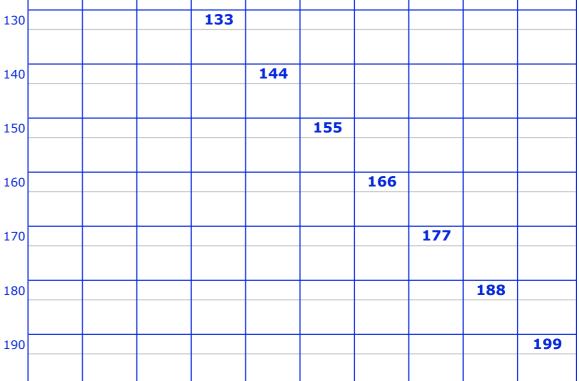
;in memory location 0,0 put 0
;in memory location 0,1 put 100
;in memory location 0,2 put 1
;in memory location 0,3 put 10
; load R1 with contents of

memory 00

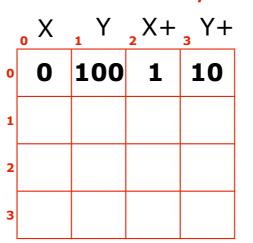
; load R2 with contents of memory 01

;add and put result in R3

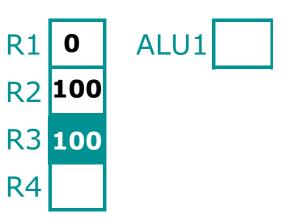




#### Main Memory



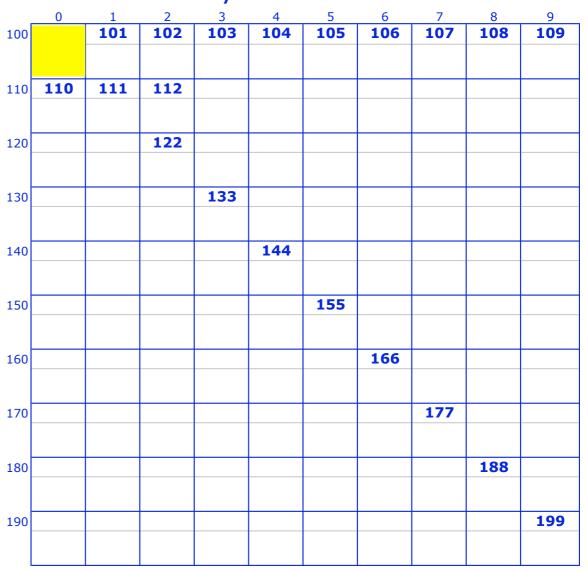
#### **Microprocessor**



;in memory location 0,0 put 0
;in memory location 0,1 put 100
;in memory location 0,2 put 1
;in memory location 0,3 put 10
; load R1 with contents of
memory 00
; load R2 with contents of
memory 01

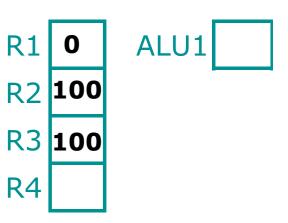
;add and put result in R3

; put 1 (to light a pixel) in the address in R3



## 

#### Microprocessor



; in memory location 0,0 put 0 ; in memory location 0,1 put 100

; in memory location 0,2 put 1

; in memory location 0,3 put 10

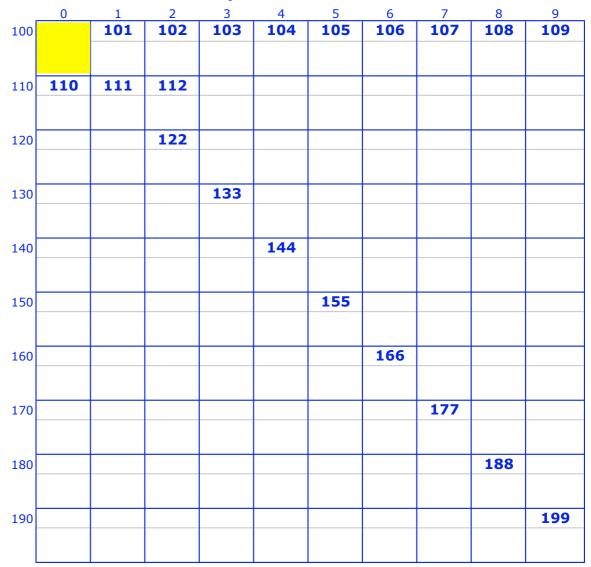
; load R1 with contents of memory 00

; load R2 with contents of memory 01

;add and put result in R3

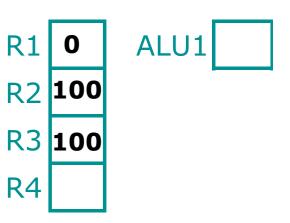
; put 1 (to light a pixel) in the address in R3

; increment  $\boldsymbol{x}$  with contents of memory 02



## 

#### Microprocessor



; in memory location 0,0 put 0 ; in memory location 0,1 put 100

; in memory location 0,2 put 1

; in memory location 0,3 put 10

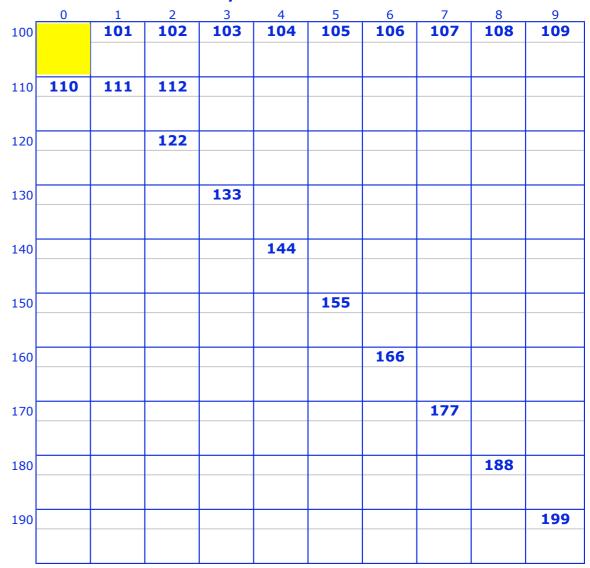
; load R1 with contents of memory 00

; load R2 with contents of memory 01

;add and put result in R3

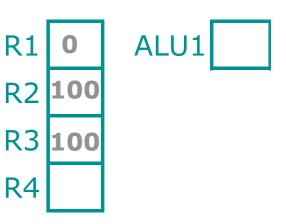
; put 1 (to light a pixel) in the address in R3

; increment  $\boldsymbol{x}$  with contents of memory 02



## Main Memory <sup>0</sup> X <sup>1</sup> Y <sup>2</sup> X+ <sup>3</sup> Y+ **0 100 1 10** <sup>1</sup> 10 10 <sup>1</sup> 10 10 <sup>2</sup> 10 10

#### Microprocessor



;in memory location 0,0 put 0 ;in memory location 0,1 put 100

; in memory location 0,2 put 1

; in memory location 0,3 put 10

; load R1 with contents of memory 00

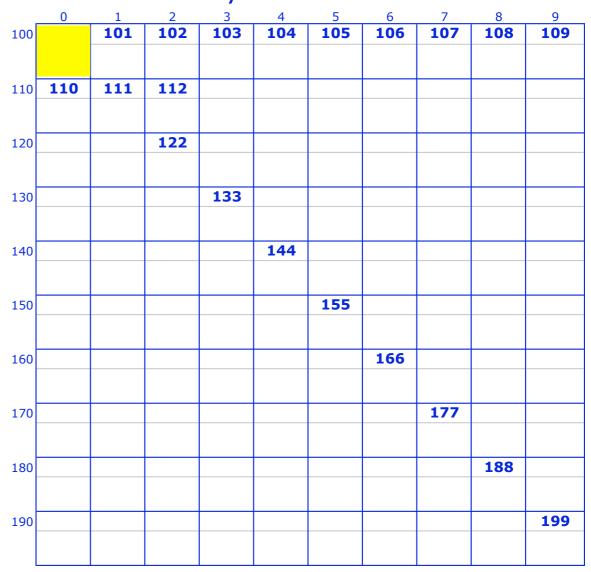
; load R2 with contents of memory 01

;add and put result in R3

; put 1 (to light a pixel) in the address in R3

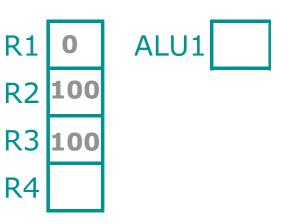
; increment x with contents of memory 02

; increment y with contents of memory 03



## Main Memory <sup>0</sup> X <sup>1</sup> Y <sup>2</sup> X+<sup>3</sup> Y+ **0 110 1 10** <sup>1</sup> 1 10 <sup>2</sup> 1 1 10 <sup>3</sup> 1 1 10

#### Microprocessor



;in memory location 0,0 put 0 ;in memory location 0,1 put 100

; in memory location 0,2 put 1

; in memory location 0,3 put 10

; load R1 with contents of memory 00

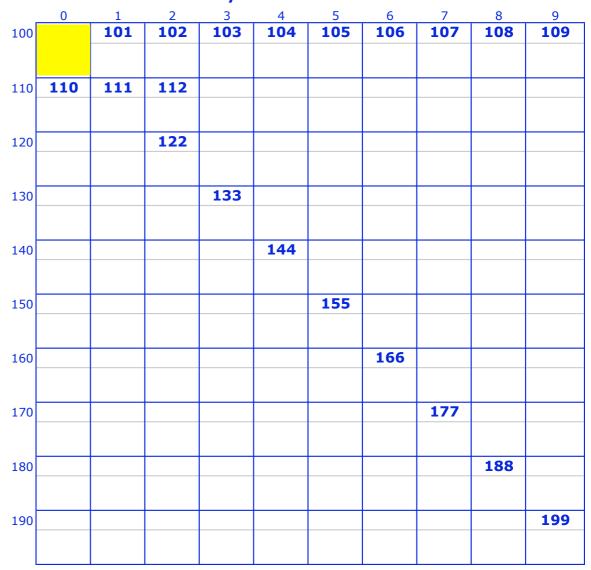
; load R2 with contents of memory 01

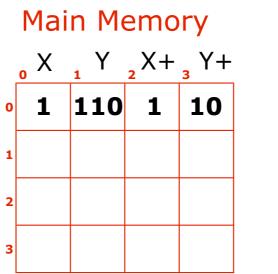
;add and put result in R3

; put 1 (to light a pixel) in the address in R3

; increment x with contents of memory 02

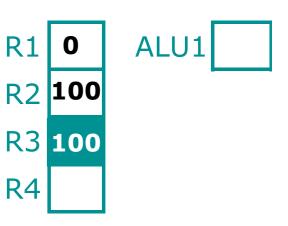
; increment y with contents of memory 03





190

#### Microprocessor



; in memory location 0,0 put 0

; in memory location 0,1 put 100

; in memory location 0,2 put 1

; in memory location 0,3 put 10

; load R1 with contents of memory 00

; load R2 with contents of memory 01

;add and put result in R3

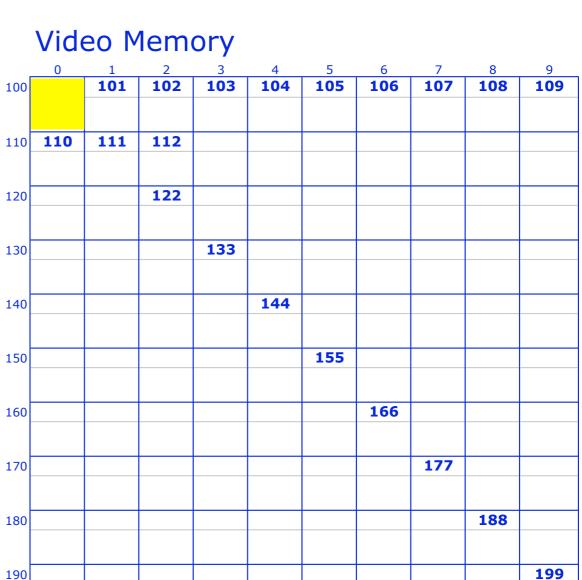
; put 1 (to light a pixel) in the address in R3

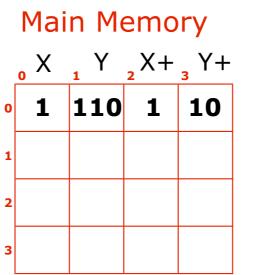
; increment x with contents of memory 02

; increment y with contents of memory 03

; check if R3=199 (R3 is where the video memory location was stored);

; jump to label 'DRAW'



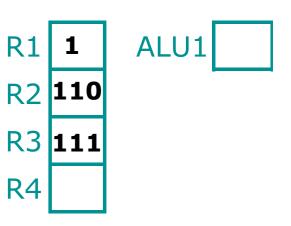


170

180

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#### Microprocessor



177

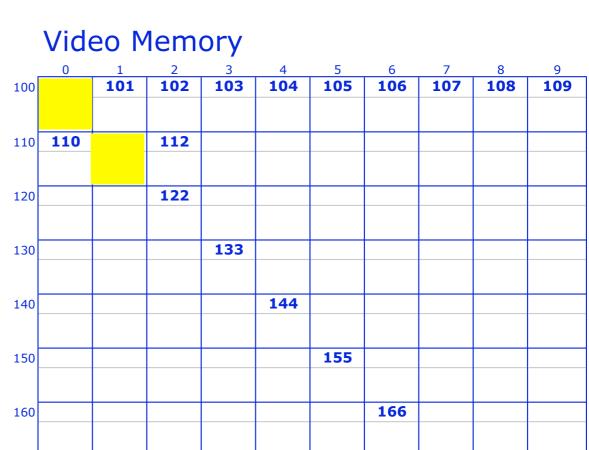
188

199

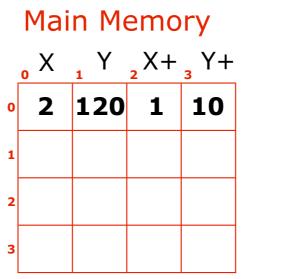
;in memory location 0,0 put 0
;in memory location 0,1 put 100
;in memory location 0,2 put 1
;in memory location 0,3 put 10
; load R1 with contents of
memory 00

; load R2 with contents of memory 01 ; add and put result in R3

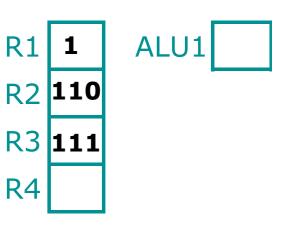
; put 1 (to light a pixel) in the address in R3



DRAW



#### Microprocessor



;in memory location 0,0 put 0
;in memory location 0,1 put 100
;in memory location 0,2 put 1
;in memory location 0,3 put 10

DRAW

; load R1 with contents of memory 00

; load R2 with contents of memory 01

;add and put result in R3

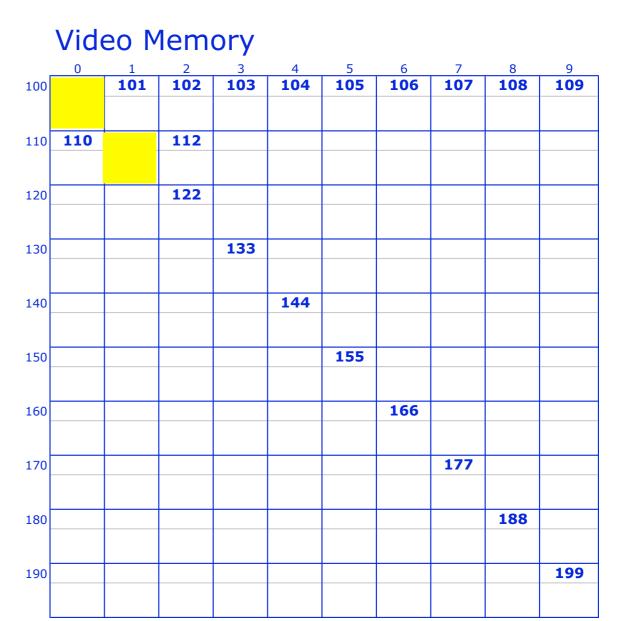
; put 1 (to light a pixel) in the address in R3

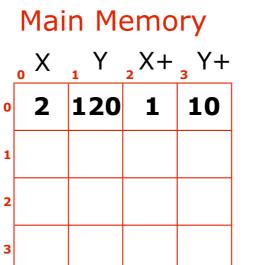
; increment  $\boldsymbol{x}$  with contents of memory 02

; increment y with contents of memory 03

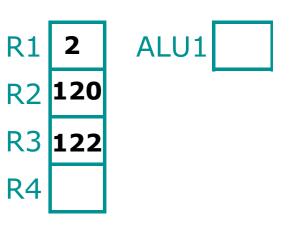
; check if R3=199 (R3 is where the video memory location was stored);

; jump to label 'DRAW'

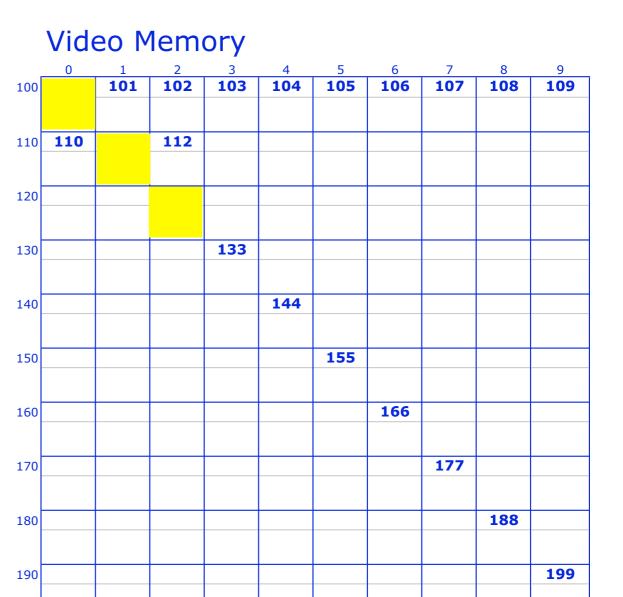




#### Microprocessor



;in memory location 0,0 put 0
;in memory location 0,1 put 100
;in memory location 0,2 put 1
;in memory location 0,3 put 10
; load R1 with contents of
memory 00
; load R2 with contents of
memory 01
;add and put result in R3
; put 1 (to light a pixel) in the
address in R3



#### **Terese mini-computer: example of assembler program**

Label opcode operand 1 operand 2 operand 3 ; in memory location 0,0 put 0 X: C&M 00 **#0** 1 ; in memory location 0,1 put 100 2 Y: C&M 01 #100 ; in memory location 0,2 put 1 PLUS X: C&M 3 02 #1 ; in memory location 0,3 put 10 PLUS Y: C&M 03 **#10** 4 ; load R1 with contents of memory 00 5 **DRAW:** LR1 00 ; load R2 with contents of memory 01 6 LR2 01 ;add and put result in R3 7 ADD **R1 R2 R3** ; put 1 (to light a pixel) in the address in R3 8 C&M **R3** #1 ; increment x with contents of memory 02 9 INC 02 X ; increment y with contents of memory 03 10 INC 03 Y ; check if R3=199 (R3 is where the video memory location was stored); END JEQ 11 **#R3** 199 ; jump to label 'DRAW' 12 JMP **DRAW** 13 END: END

